ECO316: Applied game theory Lecture 11

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low price	10,0	3,3

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- After each play of game, each participant will be informed of his/her opponent's action
- Number of periods determined randomly
- After first 10 periods, game will end with probability 0.07 after each period

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 - players may be "impatient" (prefer to consume now rather than later)
 - game may end with positive probability after any period
- Model allows us to think about long-term relationships and reputation

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where $0 < \delta < 1$

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Example: suppose outcome in stage game is the same in every period: a, a, a, ...

- Assume players discount future payoffs
- > Player *i*'s payoff function in stage game: u_i
- ▶ Player *i*'s payoff to sequence a^1, a^2, a^3, \ldots of action profiles in stage game:

$$u_i(a^1) + \delta u_i(a^2) + \delta^2 u_i(a^3) + \cdots$$

where $0 < \delta < 1$

Example: suppose outcome in stage game is the same in every period: a, a, a, ...

Then payoff in repeated game is

 $u_i(a) + \delta u_i(a) + \delta^2 u_i(a) + \cdots =$

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- Player i's payoff function in stage game: u_i
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$$u_i(a) + \delta u_i(a) + \delta^2 u_i(a) + \cdots = (1 + \delta + \delta^2 + \cdots) u_i(a)$$

- Assume players discount future payoffs
- Player i's payoff function in stage game: u_i
- Player i's payoff to sequence a¹, a², a³,... of action profiles in stage game:

$$u_i(a^1) + \delta u_i(a^2) + \delta^2 u_i(a^3) + \cdots$$

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$$egin{aligned} u_i(m{a}) + \delta u_i(m{a}) + \delta^2 u_i(m{a}) + \cdots &= (\mathbf{1} + \delta + \delta^2 + \cdots) u_i(m{a}) \ &= rac{1}{\mathbf{1} - \delta} u_i(m{a}) \end{aligned}$$

Strategy in repeated game is function that specifies

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- action in third period for all possible action profiles in first and second periods

Strategy in repeated game is function that specifies

 action in each period for every possible history of action profiles

- action in first period
- action in second period for every possible action profile in first period
- action in third period for all possible action profiles in first and second periods
- and so on ...

Repeated games: Strategies Example: Stage game is *Prisoner's Dilemma*

	С	D
С	2,2	0,3
D	3,0	1,1

Some possible strategies in repeated game:
$$\begin{array}{c|c}
C & D \\
C & 2,2 & 0,3 \\
D & 3,0 & 1,1
\end{array}$$

Some possible strategies in repeated game:

All-C Choose *C* in the first period, and in every future period, for every possible history

$$\begin{array}{c|cc}
C & D \\
\hline
C & 2,2 & 0,3 \\
\hline
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\end{array}$$

Some possible strategies in repeated game:

- *All-C* Choose *C* in the first period, and in every future period, for every possible history
- *All-D* Choose *D* in the first period, and in every future period, for every possible history

$$\begin{array}{c|cc}
C & D \\
\hline
C & 2,2 & 0,3 \\
\hline
D & 3,0 & 1,1 \\
\end{array}$$

Some possible strategies in repeated game:

- *All-C* Choose C in the first period, and in every future period, for every possible history
- *All-D* Choose *D* in the first period, and in every future period, for every possible history
- *C-D alternation* Choose *C* in the first period, and then alternate between *D* and *C*, regardless of the action profiles previously chosen

	С	D
С	2,2	0,3
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Some possible strategies in repeated game:

$$\begin{array}{c|c}
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Some possible strategies in repeated game:

Unrelenting punishment Choose *C* in first period, and in every future period *t* choose

C if other player chose C in every period $1, \ldots, t-1$

$$\begin{array}{c|c}
C & D \\
C & 2,2 & 0,3 \\
D & 3,0 & 1,1
\end{array}$$

Some possible strategies in repeated game:

Unrelepting punishment Choose C in first period, and in every future period t choose

- if other player chose *C* in every period $1, \ldots, t-1$ if other player chose *D* in any period $1, \ldots, t-1$

This strategy indefinitely "punishes" a player who deviates from C by choosing D subsequently

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Some possible strategies in repeated game:

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Some possible strategies in repeated game:

k-period punishment Deviation to *D* induces *k* periods of punishment

Two "states": collude, punish

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Some possible strategies in repeated game:

- Two "states": *collude*, *punish*
- Choose C in state collude, D in state punish

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C & D \\
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Some possible strategies in repeated game:

- Two "states": *collude*, *punish*
- Choose C in state collude, D in state punish
- In period 1, state is *collude*

$$\begin{array}{c|cc}
C & D \\
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Some possible strategies in repeated game:

- Two "states": *collude*, *punish*
- Choose C in state collude, D in state punish
- In period 1, state is collude
- Switch from collude to punish if opponent chooses D

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Some possible strategies in repeated game:

- Two "states": *collude*, *punish*
- Choose C in state collude, D in state punish
- In period 1, state is collude
- Switch from collude to punish if opponent chooses D
- Switch from *punish* to *collude* after *k* consecutive periods in which state is *punish*

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Some possible strategies in repeated game:

k-period punishment Deviation to *D* induces *k* periods of punishment

- Two "states": *collude*, *punish*
- Choose C in state collude, D in state punish
- In period 1, state is collude
- Switch from collude to punish if opponent chooses D
- Switch from *punish* to *collude* after *k* consecutive periods in which state is *punish*

Strategy "punishes" a player who deviates from C by playing D for k periods

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Some possible strategies in repeated game:

Tit-for-tat

	С	D
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Some possible strategies in repeated game:

Tit-for-tat Choose *C* in first period, and in every future period choose action chosen by other player in previous period

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Axelrod's	experir	ments		

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- Axelrod wrote book The evolution of cooperation in which he argued the evolutionary merits of tit-for-tat
- Despite *tit-for-tat*'s success and Axelrod's arguments, theoretical results showing that *tit-for-tat* is robustly good are few

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С	2,2	0,3
D	3,0	1,1

Is strategy pair (All-C, All-C) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?

	С	D
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- Is strategy pair (All-C, All-C) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-C, player 1 uses All-C \Rightarrow outcome

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 (C, C) in every period

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$$\mathsf{payoff} = \mathbf{2} + \mathbf{2}\delta + \mathbf{2}\delta^{\mathbf{2}} + \cdots =$$
	С	D
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- Is strategy pair (All-C, All-C) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
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$$\mathsf{payoff} = \mathbf{2} + 2\delta + 2\delta^2 + \dots = \frac{2}{1-\delta}$$

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- Is strategy pair (All-C, All-C) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- ► Given player 2 uses All-C, player 1 uses All-C ⇒ outcome (C, C) in every period ⇒

payoff =
$$2 + 2\delta + 2\delta^2 + \cdots = \frac{2}{1 - \delta}$$

• Player 1 deviates to $All-D \Rightarrow$ outcome

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$$\mathsf{payoff} = \mathbf{3} + \mathbf{3}\delta + \mathbf{3}\delta^2 + \cdots =$$

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$$\mathsf{payoff} = \mathbf{3} + \mathbf{3}\delta + \mathbf{3}\delta^2 + \dots = \frac{\mathbf{3}}{\mathbf{1} - \delta}$$

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- Is strategy pair (All-C, All-C) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- ► Given player 2 uses All-C, player 1 uses All-C ⇒ outcome (C, C) in every period ⇒

$$\mathsf{payoff} = 2 + 2\delta + 2\delta^2 + \dots = \frac{2}{1 - \delta}$$

$$\mathsf{payoff} = \mathbf{3} + \mathbf{3}\delta + \mathbf{3}\delta^2 + \dots = \frac{\mathbf{3}}{\mathbf{1} - \delta} > \frac{\mathbf{2}}{\mathbf{1} - \delta}$$

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- Is strategy pair (All-C, All-C) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- ► Given player 2 uses All-C, player 1 uses All-C ⇒ outcome (C, C) in every period ⇒

$$\mathsf{payoff} = 2 + 2\delta + 2\delta^2 + \dots = \frac{2}{1 - \delta}$$

► Player 1 deviates to All-D ⇒ outcome (D, C) in every period ⇒

$$\mathsf{payoff} = \mathbf{3} + \mathbf{3}\delta + \mathbf{3}\delta^2 + \dots = \frac{\mathbf{3}}{\mathbf{1} - \delta} > \frac{\mathbf{2}}{\mathbf{1} - \delta}$$

▶ So (All-C, All-C) is not a Nash equilibrium

	С	D
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Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?

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- Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-D, player 1 uses All-D ⇒ outcome
 (D, D) in every period ⇒

payoff =

	С	D
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- Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-D, player 1 uses All-D ⇒ outcome
 (D, D) in every period ⇒

payoff
$$= 1 + \delta + \delta^2 + \dots = rac{1}{1 - \delta}$$

	С	D
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- Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-D, player 1 uses All-D ⇒ outcome
 (D, D) in every period ⇒

payoff
$$= 1 + \delta + \delta^2 + \dots = \frac{1}{1 - \delta}$$

▶ Player 1 deviates to any other strategy ⇒ outcome

	С	D
С	2,2	0,3
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- Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-D, player 1 uses All-D ⇒ outcome
 (D, D) in every period ⇒

payoff
$$= 1 + \delta + \delta^2 + \dots = \frac{1}{1 - \delta}$$

► Player 1 deviates to any other strategy ⇒ outcome (C, D) or (D, D) every period ⇒ payoff

	С	D
С	2,2	0,3
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- Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-D, player 1 uses All-D ⇒ outcome
 (D, D) in every period ⇒

payoff
$$= 1 + \delta + \delta^2 + \dots = \frac{1}{1 - \delta}$$

► Player 1 deviates to any other strategy ⇒ outcome (C, D) or (D, D) every period ⇒ payoff either 0 or 1 in every period ⇒

	С	D
С	2,2	0,3
D	3,0	1,1

- Is strategy pair (All-D, All-D) a Nash equilibrium of infinitely repeated Prisoner's Dilemma?
- Given player 2 uses All-D, player 1 uses All-D ⇒ outcome
 (D, D) in every period ⇒

payoff
$$= 1 + \delta + \delta^2 + \dots = \frac{1}{1 - \delta}$$

► Player 1 deviates to any other strategy ⇒ outcome (C, D) or (D, D) every period ⇒ payoff either 0 or 1 in every period ⇒

$$\mathsf{payoff} \leq \mathsf{1} + \delta + \delta^2 + \dots = \frac{\mathsf{1}}{\mathsf{1} - \delta}$$

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So (All-D, All-D) is a Nash equilibrium

	С	D
С	2,2	0,3
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Does repeated game have Nash equilibrium that generates better outcome than (D, D) in every period?

	С	D
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- Does repeated game have Nash equilibrium that generates better outcome than (D, D) in every period?
- Suppose player 2 uses Unrelenting punishment strategy: Choose C in first period, and in every future period t choose
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- What is a best response of player 1?

Prisoner's Dilemma

	С	D
С	2,2	0,3
D	3,0	1,1

Is All-C a best response to Unrelenting punishment?

▶ Player 1 uses *All*-*C* \Rightarrow outcome is

Period	1	2	3	
Outcome				
Payoff of player 1				

Prisoner's Dilemma

Model

	С	D
С	2,2	0,3
D	3,0	1,1

Is All-C a best response to Unrelenting punishment?

► Player 1 uses All-C ⇒ outcome is (C, C) in every period ⇒ payoff

Period	1	2	3	
Outcome	(<i>C</i> , <i>C</i>)	(<i>C</i> , <i>C</i>)	(<i>C</i> , <i>C</i>)	
Payoff of player 1				

Prisoner's Dilemma

	С	D
С	2,2	0,3
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Is All-C a best response to Unrelenting punishment?

Period	1	2	3	
Outcome		(<i>C</i> , <i>C</i>)	(<i>C</i> , <i>C</i>)	
Payoff of player 1	2	2	2	

Prisoner's Dilemma

	С	D
С	2,2	0,3
D	3,0	1,1

Is All-C a best response to Unrelenting punishment?

► Player 1 uses All-C ⇒ outcome is (C, C) in every period ⇒ payoff 2 in every period

Period	1	2	3	
Outcome	(<i>C</i> , <i>C</i>)	(<i>C</i> , <i>C</i>)	(<i>C</i> , <i>C</i>)	
Payoff of player 1	2	2	2	

Any strategy that generates a different outcome chooses D after some history ((C, C), ..., (C, C))

Prisoner's Dilemma

Model

	C	D
С	2,2	0,3
D	3,0	1,1

Is All-C a best response to Unrelenting punishment?

▶ Player 1 uses All-C \Rightarrow outcome is (C, C) in every period \Rightarrow payoff 2 in every period

- Any strategy that generates a different outcome chooses D after some history $((C, C), \ldots, (C, C))$
- Say player 1 chooses C up to t 1, then D in period $t \Rightarrow$

Prisoner's Dilemma

	С	D
С	2,2	0,3
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Is All-C a best response to Unrelenting punishment?

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Prisoner's Dilemma

Model

	С	D
С	2,2	0 , 3
D	3,0	1 , 1

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Prisoner's Dilemma

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Is All-C a best response to Unrelenting punishment?

▶ Player 1 uses *All-C* ⇒ payoff in repeated game

$$= \mathbf{2} + \mathbf{2}\delta + \mathbf{2}\delta^2 + \dots$$

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▶ Player 1 chooses *C* up to period t - 1, then *D* in period $t \Rightarrow$ payoff in repeated game

$$\leq \mathbf{2} + \mathbf{2}\delta + \mathbf{2}\delta^2 + \dots + \mathbf{2}\delta^{t-2}$$

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Player 1 chooses C up to period t − 1, then D in period t ⇒ payoff in repeated game

$$\leq \mathbf{2} + \mathbf{2}\delta + \mathbf{2}\delta^2 + \dots + \mathbf{2}\delta^{t-2} + \mathbf{3}\delta^{t-1}$$

Is All-C a best response to Unrelenting punishment?

▶ Player 1 uses *All-C* ⇒ payoff in repeated game

$$= \mathbf{2} + \mathbf{2}\delta + \mathbf{2}\delta^2 + \dots$$

► Player 1 chooses C up to period t – 1, then D in period t ⇒ payoff in repeated game

$$\leq 2+2\delta+2\delta^2+\cdots+2\delta^{t-2}+3\delta^{t-1}+\delta^t+\delta^{t+1}+\cdots$$

Is All-C a best response to Unrelenting punishment?

▶ Player 1 uses *All-C* ⇒ payoff in repeated game

$$= \mathbf{2} + \mathbf{2}\delta + \mathbf{2}\delta^2 + \dots$$

Player 1 chooses C up to period t − 1, then D in period t ⇒ payoff in repeated game

$$\leq 2+2\delta+2\delta^2+\cdots+2\delta^{t-2}+3\delta^{t-1}+\delta^t+\delta^{t+1}+\cdots$$

So deviation is not profitable if

$$2 + 2\delta + \dots + 2\delta^{t-2} + 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots$$

$$\leq 2 + 2\delta + \dots + 2\delta^{t-2} + 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

Is All-C a best response to Unrelenting punishment?

$$2 + 2\delta + \dots + 2\delta^{t-2} + 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots$$

$$\leq 2 + 2\delta + \dots + 2\delta^{t-2} + 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

 \Leftrightarrow
Is All-C a best response to Unrelenting punishment?

$$2 + 2\delta + \dots + 2\delta^{t-2} + 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots$$

$$\leq 2 + 2\delta + \dots + 2\delta^{t-2} + 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

 \Leftrightarrow

$$3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots \leq 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

$$2 + 2\delta + \dots + 2\delta^{t-2} + 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots$$

$$\leq 2 + 2\delta + \dots + 2\delta^{t-2} + 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

$$\Leftrightarrow$$

$$\begin{aligned} & 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots \leq 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots \\ \Leftrightarrow \text{(divide by } \delta^{t-1}\text{)} \end{aligned}$$

Is All-C a best response to Unrelenting punishment?

$$2 + 2\delta + \dots + 2\delta^{t-2} + 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots$$

$$\leq 2 + 2\delta + \dots + 2\delta^{t-2} + 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

$$\begin{array}{l} \Leftrightarrow \\ 3\delta^{t-1} + \delta^t + \delta^{t+1} + \cdots \leq 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots \\ \Leftrightarrow \text{(divide by } \delta^{t-1}\text{)} \end{array}$$

$$3+\delta+\delta^2+\cdots\leq 2+2\delta+2\delta^2+\ldots$$

 \Leftrightarrow

Is All-C a best response to Unrelenting punishment?

$$2 + 2\delta + \dots + 2\delta^{t-2} + 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots$$

$$\leq 2 + 2\delta + \dots + 2\delta^{t-2} + 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots$$

$$\begin{array}{l} \Leftrightarrow \\ & 3\delta^{t-1} + \delta^t + \delta^{t+1} + \dots \leq 2\delta^{t-1} + 2\delta^t + 2\delta^{t+1} + \dots \\ \Leftrightarrow \text{(divide by } \delta^{t-1}\text{)} \end{array}$$

$$3+\delta+\delta^2+\cdots\leq 2+2\delta+2\delta^2+\ldots$$

 \Leftrightarrow

$$1 \le \delta + \delta^2 + \dots \quad \Leftrightarrow \quad 1 \le \frac{\delta}{1 - \delta} \quad \Leftrightarrow \quad \delta \ge \frac{1}{2}$$

$$\begin{array}{c|c}
C & D \\
\hline
C & 2,2 & 0,3 \\
\hline
D & 3,0 & 1,1
\end{array}$$

Is All-C a best response to Unrelenting punishment?

► Conclusion: if $\delta \ge \frac{1}{2}$ then the strategy *All-C* is a best response to the strategy *Unrelenting punishment*

$$\begin{array}{c|cc}
C & D \\
\hline
C & 2,2 & 0,3 \\
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D & 3,0 & 1,1 \\
\end{array}$$

- ► Conclusion: if $\delta \ge \frac{1}{2}$ then the strategy *All-C* is a best response to the strategy *Unrelenting punishment*
- $\delta \geq \frac{1}{2}$ means the players are sufficiently patient

 $\begin{array}{c|ccc}
C & D \\
C & 2, 2 & 0, 3 \\
D & 3, 0 & 1, 1
\end{array}$

- ► Conclusion: if $\delta \ge \frac{1}{2}$ then the strategy *All-C* is a best response to the strategy *Unrelenting punishment*
- $\delta \geq \frac{1}{2}$ means the players are sufficiently patient
- Is (All-C, Unrelenting punishment) a Nash equilibrium in this case?

Prisoner's Dilemma

	С	D
С	2 , 2	0 , 3
D	3,0	1,1

- ► Conclusion: if $\delta \ge \frac{1}{2}$ then the strategy *All-C* is a best response to the strategy *Unrelenting punishment*
- $\delta \geq \frac{1}{2}$ means the players are sufficiently patient
- Is (All-C, Unrelenting punishment) a Nash equilibrium in this case?
- Is Unrelenting punishment a best response to All-C?

Prisoner's Dilemma

	С	D
С	2 , 2	0, <mark>3</mark>
D	3,0	1,1

- ► Conclusion: if $\delta \ge \frac{1}{2}$ then the strategy *All-C* is a best response to the strategy *Unrelenting punishment*
- $\delta \geq \frac{1}{2}$ means the players are sufficiently patient
- Is (All-C, Unrelenting punishment) a Nash equilibrium in this case?
- ▶ Is Unrelenting punishment a best response to All-C?
- No! If player 1 chooses C regardless of player 2's actions, player 2 is better off choosing D in every period

Prisoner's Dilemma

	С	D
С	2 , 2	0 , 3
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- ► Conclusion: if $\delta \ge \frac{1}{2}$ then the strategy *All-C* is a best response to the strategy *Unrelenting punishment*
- $\delta \geq \frac{1}{2}$ means the players are sufficiently patient
- Is (All-C, Unrelenting punishment) a Nash equilibrium in this case?
- Is Unrelenting punishment a best response to All-C?
- No! If player 1 chooses C regardless of player 2's actions, player 2 is better off choosing D in every period
- So (All-C, Unrelenting punishment) is not a Nash equilibrium of the repeated game

Prisoner's Dilemma



Player 1's payoffs in repeated game

- Best responses to Unrelenting punishment
 - All-C is not the only best response to Unrelenting punishment

Prisoner's Dilemma

Stage game



Player 1's payoffs in repeated game

Best responses to Unrelenting punishment

- All-C is not the only best response to Unrelenting punishment
- Any strategy of player 1 that generates the outcome path (C, C), (C, C), ... against Unrelenting punishment is also a best response

Prisoner's Dilemma

Stage game



Player 1's payoffs in repeated game

Best responses to Unrelenting punishment

- All-C is not the only best response to Unrelenting punishment
- Any strategy of player 1 that generates the outcome path (C, C), (C, C), ... against Unrelenting punishment is also a best response
- Unrelenting punishment is such a strategy

Prisoner's Dilemma

Stage game



Player 1's payoffs in repeated game

Best responses to Unrelenting punishment

- All-C is not the only best response to Unrelenting punishment
- Any strategy of player 1 that generates the outcome path (C, C), (C, C), ... against Unrelenting punishment is also a best response
- Unrelenting punishment is such a strategy
- ► Thus (Unrelenting punishment, Unrelenting punishment) is a Nash equilibrium of the repeated game if $\delta \ge \frac{1}{2}$

	С	D
С	2,2	0,3
D	3,0	1,1

- Unrelenting punishment: Choose C in first period, and in every future period t choose
 - if other player chose *C* in every period $1, \ldots, t-1$ if other player chose *D* in any period $1, \ldots, t-1$

	С	D
С	2,2	0,3
D	3,0	1,1

- Unrelenting punishment: Choose C in first period, and in every future period t choose
 - $\begin{cases} C & \text{if other player chose } C \text{ in every period } 1, \dots, t-1 \\ D & \text{if other player chose } D \text{ in any period } 1, \dots, t-1 \end{cases}$
- If $\delta \geq \frac{1}{2}$ then strategy pair (Unrelenting punishment, Unrelenting punishment) is Nash equilibrium of infinitely repeated Prisoner's Dilemma with payoffs shown above

	С	D
С	2,2	0,3
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- If $\delta \geq \frac{1}{2}$ then strategy pair (Unrelenting punishment, Unrelenting punishment) is Nash equilibrium of infinitely repeated Prisoner's Dilemma with payoffs shown above
- ⇒ "collusion" possible if players sufficiently patient

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С	2,2	0,3
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- If $\delta \geq \frac{1}{2}$ then strategy pair (*Unrelenting punishment*, Unrelenting punishment) is Nash equilibrium of infinitely repeated Prisoner's Dilemma with payoffs shown above
- \Rightarrow "collusion" possible if players sufficiently patient
 - For Prisoner's Dilemma with different payoffs, cutoff δ differs

	С	D
С	2,2	0,3
D	3,0	1,1

Less severe punishment?

Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?

	С	D
С	2,2	0,3
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- Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?
- Suppose player 2 uses k-period punishment:

	С	D
С	2,2	0,3
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- Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?
- Suppose player 2 uses *k*-period punishment:
 - Choose C in first period

	С	D
С	2,2	0,3
D	3,0	1,1

- Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?
- Suppose player 2 uses k-period punishment:
 - Choose C in first period
 - Continue to choose C until player 1 chooses D

	С	D
С	2,2	0,3
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- Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?
- Suppose player 2 uses k-period punishment:
 - Choose C in first period
 - Continue to choose C until player 1 chooses D
 - Starting in period after player 1 chooses D, choose D for k periods (regardless of player 1's actions)

	С	D
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- Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?
- Suppose player 2 uses k-period punishment:
 - Choose C in first period
 - Continue to choose C until player 1 chooses D
 - Starting in period after player 1 chooses D, choose D for k periods (regardless of player 1's actions)
 - After k periods of choosing D, choose C, but again start a k-period punishment (choosing D) if player 1 chooses D in any period



- Does repeated game have Nash equilibrium with outcome (C, C) in every period in which punishment is less severe?
- Suppose player 2 uses *k*-period punishment:
 - Choose C in first period
 - Continue to choose C until player 1 chooses D
 - Starting in period after player 1 chooses D, choose D for k periods (regardless of player 1's actions)
 - After k periods of choosing D, choose C, but again start a k-period punishment (choosing D) if player 1 chooses D in any period
- Does this strategy deter deviations by player 1 from C?

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	 t + k	<i>t</i> + <i>k</i> + 1
Outcome	(<i>C</i> , <i>C</i>)	 (<i>C</i> , <i>C</i>)					
Payoff of player 1	2	 2					

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	 t + k	t + k + 1
Outcome	(<i>C</i> , <i>C</i>)	 (<i>C</i> , <i>C</i>)	(D, C)	(?, D)	(?, D)	 (?, D)	
Payoff of player 1	2	 2	3	\leq 1	\leq 1	 \leq 1	

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	 t + k	<i>t</i> + <i>k</i> + 1
Outcome	(<i>C</i> , <i>C</i>)	 (<i>C</i> , <i>C</i>)	(D, C)	(?, D)	(?, D)	 (?, D)	(?, C)
Payoff of player 1	2	 2	3	\leq 1	\leq 1	 \leq 1	?

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	 t + k	<i>t</i> + <i>k</i> + 1
Outcome							
Payoff of player 1	2	 2	3	\leq 1	\leq 1	 \leq 1	?

If player 1 chooses C in period t + k + 1 and subsequently then outcome from then on is (C, C)

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	 t + k	<i>t</i> + <i>k</i> + 1
Outcome	(<i>C</i> , <i>C</i>)	 (<i>C</i> , <i>C</i>)	(D, C)	(?, D)	(?, D)	 (?, D)	(?, C)
Payoff of player 1	2	 2	3	\leq 1	\leq 1	 \leq 1	?

- If player 1 chooses C in period t + k + 1 and subsequently then outcome from then on is (C, C)
- If player 1 chooses D in period t + k + 1 then new cycle of punishment starts

Less severe punishment?

▶ If player 1 deviates to *D* in period *t*, sequence of outcomes:

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	 t + k	t + k + 1
Outcome	(<i>C</i> , <i>C</i>)	 (<i>C</i> , <i>C</i>)	(<i>D</i> , <i>C</i>)	(?, D)	(?, D)	 (?, D)	(?, C)
Payoff of player 1	2	 2	3	≤ 1	≤ 1	 < 1 ≤ 1	?

- If player 1 chooses C in period t + k + 1 and subsequently then outcome from then on is (C, C)
- If player 1 chooses D in period t + k + 1 then new cycle of punishment starts
- If profitable deviation exists, then single deviation to D, with return to C, must be profitable

- Less severe punishment?
 - For deviation not to be profitable, need

$$3 + \delta + \delta^2 + \dots + \delta^k \leq 2 + 2\delta + 2\delta^2 + \dots + 2\delta^k$$

- Less severe punishment?
 - For deviation not to be profitable, need

$$\begin{aligned} \mathbf{3} + \delta + \delta^2 + \cdots + \delta^k &\leq \mathbf{2} + 2\delta + 2\delta^2 + \cdots + 2\delta^k \\ \mathbf{2} + (\mathbf{1} + \delta + \delta^2 + \cdots + \delta^k) &\leq \mathbf{2}(\mathbf{1} + \delta + \delta^2 + \cdots + \delta^k) \end{aligned}$$

Less severe punishment?

For deviation not to be profitable, need

$$\begin{aligned} \mathbf{3} + \delta + \delta^2 + \cdots + \delta^k &\leq \mathbf{2} + 2\delta + 2\delta^2 + \cdots + 2\delta^k \\ \mathbf{2} + (\mathbf{1} + \delta + \delta^2 + \cdots + \delta^k) &\leq \mathbf{2}(\mathbf{1} + \delta + \delta^2 + \cdots + \delta^k) \\ \mathbf{2} &\leq \mathbf{1} + \delta + \delta^2 + \cdots + \delta^k \end{aligned}$$

Less severe punishment?

For deviation not to be profitable, need

$$\begin{aligned} \mathbf{3} + \delta + \delta^2 + \cdots + \delta^k &\leq \mathbf{2} + 2\delta + 2\delta^2 + \cdots + 2\delta^k \\ \mathbf{2} + (\mathbf{1} + \delta + \delta^2 + \cdots + \delta^k) &\leq \mathbf{2}(\mathbf{1} + \delta + \delta^2 + \cdots + \delta^k) \\ \mathbf{2} &\leq \mathbf{1} + \delta + \delta^2 + \cdots + \delta^k \\ \mathbf{2} &\leq \frac{\mathbf{1} - \delta^{k+1}}{\mathbf{1} - \delta} \end{aligned}$$
Less severe punishment?

For deviation not to be profitable, need

$$\begin{aligned} 3+\delta+\delta^2+\dots+\delta^k&\leq 2+2\delta+2\delta^2+\dots+2\delta^k\\ 2+(1+\delta+\delta^2+\dots+\delta^k)&\leq 2(1+\delta+\delta^2+\dots+\delta^k)\\ &2&\leq 1+\delta+\delta^2+\dots+\delta^k\\ &2&\leq \frac{1-\delta^{k+1}}{1-\delta}\\ &1-2\delta+\delta^{k+1}&\leq 0\end{aligned}$$

Less severe punishment?

For deviation not to be profitable, need

$$\begin{aligned} 3+\delta+\delta^2+\dots+\delta^k&\leq 2+2\delta+2\delta^2+\dots+2\delta^k\\ 2+(1+\delta+\delta^2+\dots+\delta^k)&\leq 2(1+\delta+\delta^2+\dots+\delta^k)\\ &2&\leq 1+\delta+\delta^2+\dots+\delta^k\\ &2&\leq \frac{1-\delta^{k+1}}{1-\delta}\\ &1-2\delta+\delta^{k+1}&\leq 0\end{aligned}$$

► $k = 1 \Rightarrow 1 - 2\delta + \delta^2 \le 0 \Rightarrow (1 - \delta)^2 \le 0 \Rightarrow$ impossible!

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k = 1 ⇒ 1 − 2δ + δ² ≤ 0 ⇒ (1 − δ)² ≤ 0 ⇒ impossible! *k* = 2 ⇒ 1 − 2δ + δ³ ≤ 0 ⇒ δ ≥ 0.62

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k = 1 ⇒ 1 − 2δ + δ² ≤ 0 ⇒ (1 − δ)² ≤ 0 ⇒ impossible! *k* = 2 ⇒ 1 − 2δ + δ³ ≤ 0 ⇒ δ ≥ 0.62 *k* ↑ ⇒ cutoff value of δ ↓ ¹/₂

 $\begin{array}{c|c}
C & D \\
C & 2,2 & 0,3 \\
D & 3,0 & 1,1
\end{array}$

Less severe punishment: Conclusion

For any value of k ≥ 2, strategy pair in which each player punishes other for k periods in event of deviation is Nash equilibrium of infinitely repeated game if δ is large enough

Experiment	Model	Axelrod's experiments	Nash equilibrium	Duopoly
Repeate	ed gam	nes: Nash equi	librium	

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• Larger $k \Rightarrow$ smaller lower bound on δ

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Less severe punishment: Conclusion

- For any value of k ≥ 2, strategy pair in which each player punishes other for k periods in event of deviation is Nash equilibrium of infinitely repeated game if δ is large enough
- Larger $k \Rightarrow$ smaller lower bound on δ
- Mutually desirable outcome (C, C) is sustained by short punishment only if players are relatively patient



Summary

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- If players are sufficiently patient, also has Nash equilibrium in which outcome is (C, C) every period
 - In equilibrium strategy pair, each player "punishes" other for choosing D by playing D for some time
 - If players are more patient, punishment can be shorter

1

Repeated games: Nash equilibrium Prisoner's Dilemma

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 By similar arguments, other outcome paths can generated by Nash equilibria

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 - player 1's payoff is 0 for this path

Model

Repeated games: Nash equilibrium Prisoner's Dilemma



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- ▶ Example: (*C*, *C*), (*D*, *D*), (*C*, *C*), (*D*, *D*), ...
- Does any Nash equilibrium generate outcome path (C, D), (C, D), (C, D), ...?
- ► No:
 - player 1's payoff is 0 for this path
 - ▶ if she chooses *All-D* her payoff is 1 in every period

Experiment	Model	Axelrod's experiments	Nash equilibrium	Duopoly
Repeate	ed gam	es: Nash equi	librium	

	С	D
С	2,2	0,3
D	3,0	1,1

Other equilibrium outcome paths

 Because All-D yields player 1 payoff of at least 1 in every period, her payoff in Nash equilibrium cannot be less than 1/(1 - δ)

- Two firms
- Same constant unit cost: $C_i(q_i) = cq_i$ for i = 1, 2
- Demand function: D(p)
- Profit of firm i:

$$\pi_i(p_1, p_2) = \begin{cases} (p_i - c)D(p_i) & \text{if } p_i < p_j \\ \frac{1}{2}(p_i - c)D(p_i) & \text{if } p_i = p_j \\ 0 & \text{if } p_i > p_j \end{cases}$$

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- Firms interact repeatedly
- ▶ Each firm discounts future profits at rate δ , with $0 < \delta < 1$

Model: Price-setting firms

Denote monopoly price by p^m:

Model: Price-setting firms

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 p^m maximizes (p - c)D(p)

• Denote monopoly profit by Π^m

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Denote monopoly price by p^m:

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- Strategy:
 - In period 1 choose price p^m
 - In period *t* ≥ 2,
 - choose p^m if other firm chose p^m in every previous period
 - choose *c* if other firm chose a price different from *p^m* in some previous period

Model: Price-setting firms

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Model

Repeated games: Nash equilibrium Application: Duopoly

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- Suppose firm 2 uses the strategy
- ► If firm 1 uses the strategy, then outcome is (p^m, p^m) in every period \Rightarrow firm 1's profit is $\frac{1}{2}\Pi^m$ in every period
- If firm 1 deviates from strategy in period t, firm 2 chooses price c subsequently, regardless of firm 1's price in period t

Period	1	 <i>t</i> – 1		t	<i>t</i> + 1	<i>t</i> + 2	
Outcome	(p^{m}, p^{m})	 (p^{m}, p^{m})	(, p ^m)	(?, c)	(?, c)	
Firm 1's profit	$\frac{1}{2}\Pi^m$	 $\frac{1}{2}\Pi^m$					

Model

Repeated games: Nash equilibrium Application: Duopoly

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Firm 1's profit	$\frac{1}{2}\Pi^m$	 $\frac{1}{2}\Pi^m$				

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- If firm 1 deviates from strategy in period t, firm 2 chooses price c subsequently, regardless of firm 1's price in period t
- \Rightarrow good deviation is price slightly less than p^m
- \Rightarrow firm 1's profit is close to Π^m in period t

Period	1	 <i>t</i> – 1	t	<i>t</i> + 1	<i>t</i> + 2	
Outcome Firm 1's profit	(p^m, p^m)	 (p^m, p^m)	$(p^m - \varepsilon, p^m)$	(?, c)	(?, c)	
Firm 1's profit	$\frac{1}{2}\Pi^m$	 $\frac{1}{2}\Pi^m$	$\approx \Pi^m$	0	0	

Model: Price-setting firms

Firm 1's payoff if it uses the strategy:

$$\frac{1}{2}\Pi^m + \frac{1}{2}\delta\Pi^m + \dots + \frac{1}{2}\delta^{t-2}\Pi^m + \frac{1}{2}\delta^{t-1}\Pi^m + \frac{1}{2}\delta^t\Pi^m + \dots$$
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So deviation is not profitable if

$$\delta^{t-1}\Pi^m \leq \frac{1}{2}\delta^{t-1}\Pi^m + \frac{1}{2}\delta^t\Pi^m + \dots$$
$$= \frac{1}{2}\delta^{t-1}(\Pi^m + \delta\Pi^m + \dots) = \frac{1}{2}\delta^{t-1}\frac{\Pi^m}{1-\delta}$$

Model: Price-setting firms

So deviation by firm 1 is not profitable if

$$\delta^{t-1} \Pi^m \le \frac{1}{2} \delta^{t-1} \frac{\Pi^m}{1-\delta}$$

or

Repeated games: Nash equilibrium Application: Duopoly

Model: Price-setting firms

So deviation by firm 1 is not profitable if

$$\delta^{t-1} \Pi^m \le \frac{1}{2} \delta^{t-1} \frac{\Pi^m}{1-\delta}$$
$$1 - \delta \le \frac{1}{2}$$

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Duopoly

Repeated games: Nash equilibrium Application: Duopoly

Model: Price-setting firms

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$$\delta^{t-1} \Pi^m \le \frac{1}{2} \delta^{t-1} \frac{\Pi^m}{1-\delta}$$
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$$\delta \ge \frac{1}{2}$$

Model

Repeated games: Nash equilibrium Application: Duopoly

Conclusion for model with price-setting firms

If δ ≥ 1/2 then infinitely repeated game has Nash equilibrium in which each firm charges the monopoly price as long as the other firm does so, and otherwise charges the price c

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- That is: implicit collusion is an equilibrium if firms are sufficiently patient

Model

Repeated games: Nash equilibrium Application: Duopoly

Conclusion for model with price-setting firms

- If $\delta \ge \frac{1}{2}$ then infinitely repeated game has Nash equilibrium in which each firm charges the monopoly price as long as the other firm does so, and otherwise charges the price *c*
- That is: implicit collusion is an equilibrium if firms are sufficiently patient
- As for Prisoner's Dilemma, collusive outcome can be supported by strategies with less severe punishments

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When players repeatedly interact, there are equilibrium outcomes that are very different from the equilibrium outcomes of a one-shot interaction

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- When players repeatedly interact, there are equilibrium outcomes that are very different from the equilibrium outcomes of a one-shot interaction
- In strategy pairs that generate such outcomes, each player's action in any period depend on past actions of other players
- Many equilibria exist, with wide range of outcomes
- Some equilibria correspond to "collusive" outcomes (without any explicit collusion)